Towards a Small Buffering Delay in Adaptive Video Streaming

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Motivation

State-of-the-art rate adaptation is not suitable for low-latency dynamic streaming,

due to a lack of explicit stabilization of client buffer dynamics.

In case the client buffer is at its maximum level (the maximal buffering delay),

- interactions with TCP's flow control may lead to a *biased throughput feedback*, and result in undesirably variable and low video quality;[1]
- *ON-OFF streaming pattern* occurs, and may cause unfairness with multiple video streaming sessions.[2]

In contrast to existing solutions that focus on buffer control at near-zero buffer levels, a stabilization of buffer dynamics with a filled buffer is an open issue.

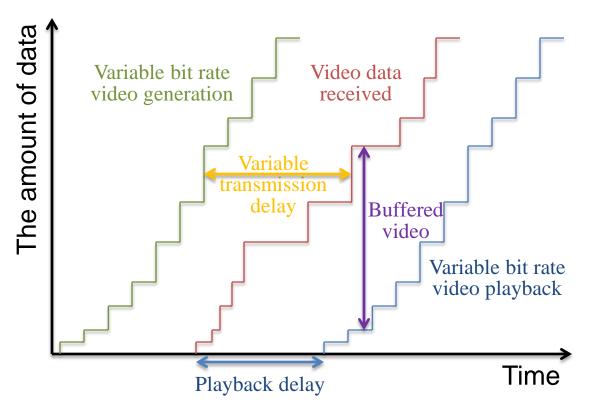
^[1] T. Huang, R. Johari, and N. McKeown. Downton abbey without the hiccups: buffer-based rate adaptation for HTTP video streaming. In *Proceedings of the ACM SIGCOMM workshop on Future human-centric multimedia networking* (FhMN), 2013.

^[2] S. Akhshabi, L. Anantakrishnan, A. Begen, and C. Dovrolis. What happens when HTTP adaptive streaming players compete for bandwidth?. In *Proceedings of the 22nd international workshop on Network and Operating System Support for Digital Audio and Video* (NOSSDAV), 2012.

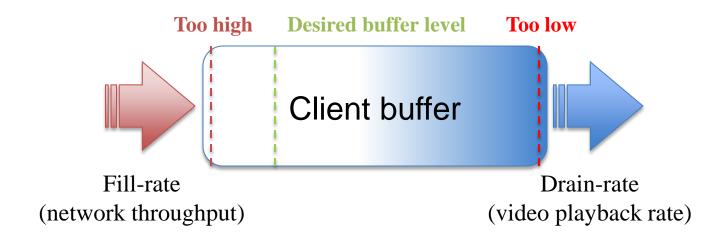
Outline

- **✓** Motivation
- ☐ Rate control for buffer stabilization
 - Buffering delay
 - Modeling buffer dynamics
 - Rate selection
- **☐** Prototype implementation
 - Server-based streaming architecture (Open-Loop rAte Control, OLAC)
 - Transport protocol configuration
- ☐ Results

Buffering Delay

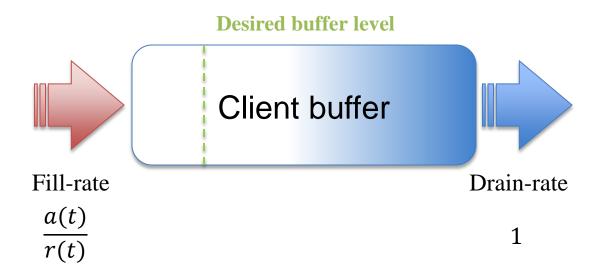


- Buffering Delay is buffered video in seconds.
- We achieve *low-latency* dynamic video streaming with buffering delays as low as the chunk-duration.

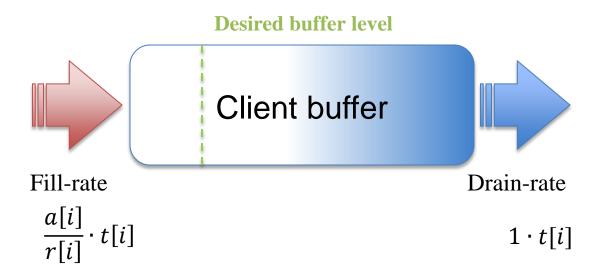


Quality Selection

Stabilizing the buffer to the desired level by regulating the drain-rate, i.e. by selecting a video bit rate for the chunk.



Express the buffer level in seconds of video.

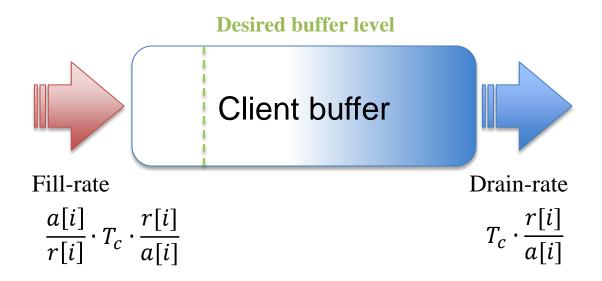


Compute the buffer level every discrete chunk.

a[i]: the throughput rate achieved during the reception of chunk i

r[i]: the selected video bit rate of chunk i

t[i]: the reception duration of chunk i

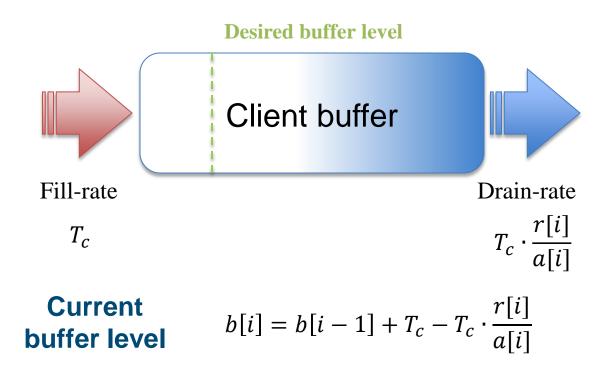


Compute the buffer level every discrete chunk.

a[i]: the throughput rate achieved during the reception of chunk i

r[i]: the selected video bit rate of chunk i

 T_c : the chunk duration



Rate Selection

Buffer dynamics

$$b_R[i] = b[i-1] + T_c - T_c \cdot \frac{r_R[i]}{a[i]}$$

Rate selection

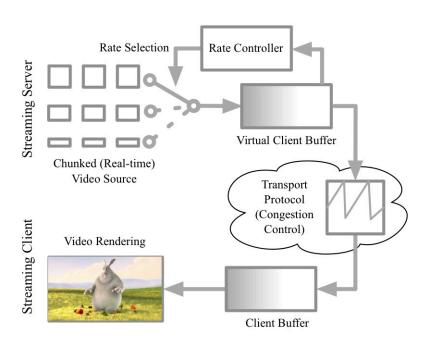
$$\widehat{R}[i] = \underset{R \in \mathcal{R}}{\operatorname{argmin}} |b_R[i] - \beta_{ref}|$$

b[i]	:	the buffer level (in seconds) when the client finishes the reception of chunk <i>i</i>	$\widehat{R}[i]$:	the selected quality level (the nominal bit rate) of the video for chunk <i>i</i>	
T_c	:	the chunk duration (each chunk containing a fixed duration of video)	R	:	the quality level (the nominal bit rate) of the video	
$r_R[i]$:	the selected video bit rate for chunk i with the nominal bit rate R	${\cal R}$:	the set of quality levels (nominal bit rates) of the video	
a[i]	:	the throughput rate achieved for chunk i	β_{ref}	:	the desired buffer level (in seconds)	11

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Open-Loop rAte Control (OLAC) [3]



- **Virtual client buffer** simulates client buffer on the server.
- A rate control on the server offers immediate feedback from clients.
- Hybrid throughput- and buffer-based adaptation balances efficiency and stability.

Transport Protocol Configurations

Our streaming prototype implementation is evaluated with two transport protocol configurations: standard **TCP-Cubic** and Predictably Reliable Real-time Transport (**PRRT**).

PRRT [4] provides

- error control under a specific delay constraint (*Predictable Reliability*),
- adaptive proactive and reactive error control (*capacity-approaching*),
- opportunistic TCP-friendliness by delay and equation-based congestion control,
- and accurate throughput estimate for applications.

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Performance Comparison

Our benchmark rate controls are

- **DASH** VLC plugin [5] and
- Quality Adaptation Controller (QAC) for adaptive video streaming [6].

We deploy the buffer stabilizer within OLAC streaming architecture on top of

- TCP, referring to Dynamic Adaptive Streaming over TCP (**DAST**), and
- PRRT, referring to Dynamic Adaptive Streaming over PRRT (**DASP**).

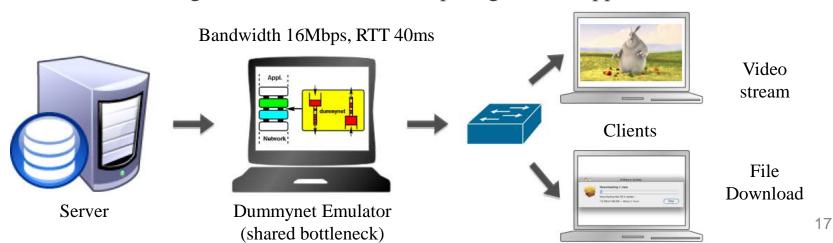
Therefore, our performance comparison contains four sets of performance results: DASH, QAC, DAST (OLAC over TCP), and DASP (OLAC over PRRT).

^[5] C. Müller and C. Timmerer. A VLC media player plugin enabling dynamic adaptive streaming over HTTP. In *Proceedings of the 19th ACM international conference on Multimedia* (MM), Scottsdale, USA, 2011.

^[6] L. Cicco, S. Mascolo, and V. Palmisano. Feedback control for adaptive live video streaming. In *Proceedings of the second annual ACM conference on Multimedia systems* (MMSys), San Jose, USA, 2011.

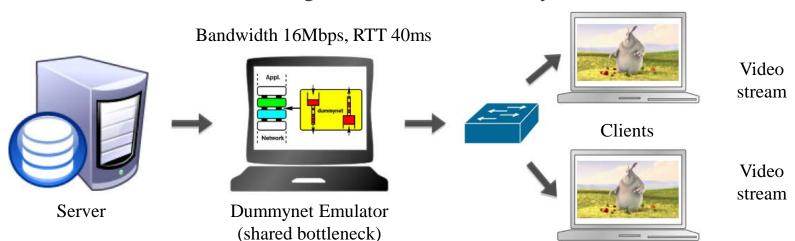
Experimental Setup

- Wide area network
- Dynamic video bit rate 1-16 Mbps, chunk duration of 2s, 4s, 6s, and 8s
- Maximum client buffer size is set to the same size of chunk duration
- Entire streaming sessions lasts 180s, competing session appears from 60-120s

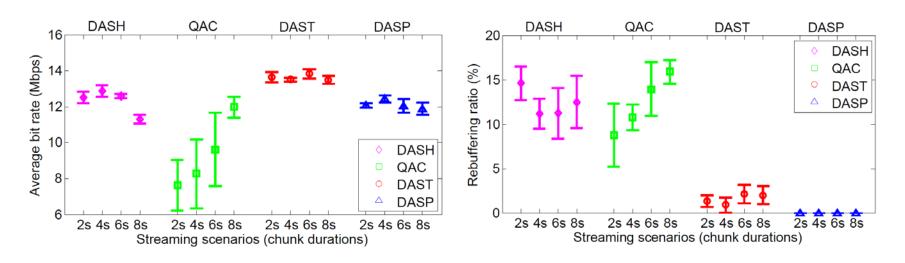


Experimental Setup

- Wide area network
- Dynamic video bit rate 1-16 Mbps, chunk duration of 2s, 4s, 6s, and 8s
- Maximum client buffer size is set to the same size of chunk duration
- Three concurrenct streaming sessions simultaneously run for 120s

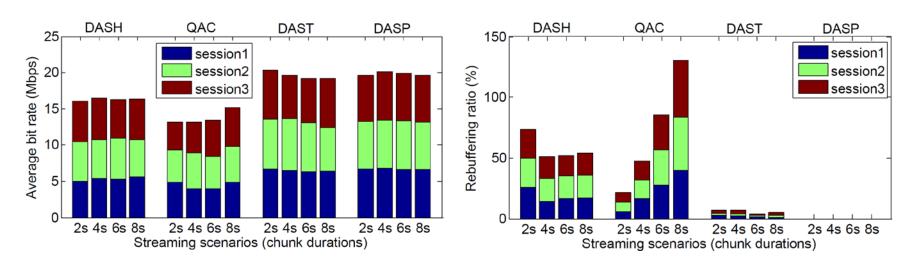


Experimental Results



- DASP had zero rebuffering events. DAST reduces the rebuffering ratio by at least 81% and 85%, compared to DASH and QAC, respectively.
- The average bit rate achieved with DAST is increased by 5-19% and 13-78% compared to DASH and QAC, respectively.

Experimental Results



- The average bit rates achieved with DAST and DASP are 17-26% and 27-54% higher compared to DASH and QAC, respectively.
- DASP had zero rebuffering events. DAST achieves with a 68-96% lower rebuffering ratio compared to QAC.

Conclusion

A solution for *low-latency* dynamic video streaming

- effectively stabilizes the buffer at a level as short as a chunk-duration,
- significantly **improves user experience** in low-latency dynamic streaming.

Thank you for your attention!